## Chapter 6: Part-2 CPU Scheduling



## Chapter 6: CPU Scheduling

- Basic Concepts
- Scheduling Criteria
- Scheduling Algorithms
- Thread Scheduling
- Multiple-Processor Scheduling
- Real-Time CPU Scheduling
- Operating Systems Examples
- Algorithm Evaluation


## Objectives

- To introduce CPU scheduling, which is the basis for multiprogrammed operating systems
- To describe various CPU-scheduling algorithms
- To discuss evaluation criteria for selecting a CPU-scheduling algorithm for a particular system
- To examine the scheduling algorithms of several operating systems


## Determining Length of Next CPU Burst

- Can only estimate the length - should be similar to the previous one
- Then pick process with shortest predicted next CPU burst
- Can be done by using the length of previous CPU bursts, using exponential averaging

1. $t_{n}=$ actual length of $n^{\text {th }} \mathrm{CPU}$ burst
2. $\tau_{n+1}=$ predicted value for the next CPU burst
3. $\alpha, 0 \leq \alpha \leq 1$
4. Define: $\quad \tau_{n=1}=\alpha t_{n}+(1-\alpha) \tau_{n}$.

- Commonly, $\alpha$ set to $1 / 2$
- Preemptive version called shortest-remaining-time-first


## Prediction of the Length of the Next CPU Burst



| CPU burst $\left(t_{i}\right)$ | 6 | 4 | 6 | 4 | 13 | 13 | 13 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| "guess" $\left(\tau_{i}\right)$ | 10 | 8 | 6 | 6 | 5 | 9 | 11 | 12 |
| $\ldots$ |  |  |  |  |  |  |  |  |

## Examples of Exponential Averaging

- $\alpha=0$
- $\tau_{n+1}=\tau_{n}$
- Recent history does not count
- $\alpha=1$
- $\tau_{n+1}=\alpha t_{n}$
- Only the actual last CPU burst counts
- If we expand the formula, we get:

$$
\begin{aligned}
\tau_{n+1}=\alpha & t_{n}+(1-\alpha) \alpha t_{n-1}+\ldots \\
& +(1-\alpha)^{\prime} \alpha t_{n-j}+\ldots \\
& +(1-\alpha)^{n+1} \tau_{0}
\end{aligned}
$$

- Since both $\alpha$ and ( $1-\alpha$ ) are less than or equal to 1 , each successive term has less weight than its predecessor


## Example of Shortest-Remaining-Time-First

- Now we add the concepts of varying arrival times and preemption to the analysis

| Process |  | Arriva/Time |  | Burst Time |
| :---: | :---: | :---: | :---: | :---: |
|  |  | 0 |  | 8 |
| $P_{1}$ |  | 1 |  |  |
| $P_{2}$ |  |  |  | 4 |
| $P_{3}$ |  | 2 |  | 9 |
| $P_{4}$ |  | 3 |  | 5 |

- Preemptive SJF Gantt Chart

| $\mathrm{P}_{1}$ | $\mathrm{P}_{2}$ | $\mathrm{P}_{4}$ | $\mathrm{P}_{1}$ | $\mathrm{P}_{3}$ |
| :--- | :--- | :--- | :--- | :--- |
| 0 | 5 | 10 | 26 |  |

- Average waiting time $=[(10-1)+(1-1)+(17-2)+5-3)] / 4=26 / 4=6.5$ msec


## Example of Shortest-Remaining-Time-First

## Another Solution :

| Process | Arrival Time | Burst Time |
| :---: | :---: | :--- |
| P1 | 0 | 876543210 |
| P2 | 1 | 43210 |
| P3 | 2 | 9876543210 |
| P4 | 3 | 543210 |



## Example of Shortest-Remaining-Time-First

Find below

- Turn-Around Time (TAT) = Complete Time (CT) - Arrival Time (AT)
- WT = Turn-Around Time (TAT) - Burst Time (BT)
- Response Time (RT) = Start Time (ST) - Arrival Time (AT)

| Process | Complete Time | Turn Around Time | Waiting Time | Response Time |
| :---: | :---: | :---: | :---: | :---: |
| P1 | 17 | $17-0=17$ | $17-8=9$ | $0-0=0$ |
| P2 | 5 | $5-1=4$ | $4-4=0$ | 1-1 =0 |
| P3 | 26 | $26-2=24$ | $24-9=15$ | $18-2=16$ |
| P4 | 10 | $10-3=7$ | $7-5=2$ | $6-3=3$ |

- Average Waiting Time $=(9+0+15+2) / 4 \rightarrow 26 / 4=6.5$
- Average Turn-Around Time (TAT) $=(17+4+24+7) / 4=13$


## Example of Shortest-Remaining-Time-First

## Another Example:

| Process | Arrival Time | Burst Time |
| :---: | :---: | :--- |
| P1 | 0 | 876543210 |
| P2 | 1 | 43210 |
| P3 | 2 | 210 |
| P4 | 3 | 10 |
| P5 | 4 | 3210 |
| P6 | 5 | 210 |


| P1 | P2 | P3 | P3 | P4 | P6 | P6 | P2 | P2 | P2 | P5 | P1 | P5 | P1 | P1 | P1 | P1 | P1 | P1 | P1 | P1 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |

## Example of Shortest-Remaining-Time-First

Find below

- Turn-Around Time (TAT) = Complete Time (CT) - Arrival Time (AT)
- WT = Turn-Around Time (TAT) - Burst Time (BT)
- Response Time (RT) = Start Time (ST) - Arrival Time (AT)

| Process | Complete Time |  | Turn Around Time |  |
| :---: | :---: | :---: | :---: | :---: |
| P1 | Waiting Time |  | Response Time |  |
| P1 | 20 | 20 | 12 | 0 |
| P2 | 10 | 9 | 5 | 0 |
| P3 | 4 | 2 | 0 | 0 |
| P4 | 5 | 2 | 1 | 1 |
| P5 | 13 | 9 | 6 | 6 |
| P6 | 7 | 2 | 0 | 0 |

- Average Waiting Time $=(12+5+0+1+6+0) / 6 \rightarrow 24 / 6=4$
- Average Turn-Around Time $($ TAT $)=(20+9+2+2+9+2) / 4 \rightarrow 44 / 4=11$


## HW of Shortest-Remaining-Time-First

- Find the average waiting time according to the SRTF (preemptive SJF) scheduling algorithm?

| Process | Arrival Time | Burst Time |
| :---: | :---: | :---: |
| P1 | 0 | 11 |
| P2 | 1 | 9 |
| P3 | 2 | 7 |
| P4 | 3 | 5 |
| P5 | 4 | 8 |

- Consider the following set of process with the length of CPU burst cycle given in milliseconds:


## HW of Shortest-Remaining-Time-First

- Find the average waiting time according to the SRTF (preemptive SJF) scheduling algorithm?

| Process | Arrival Time | Burst Time |
| :---: | :---: | :---: |
| P1 | 0 | 12 |
| P2 | 3 | 8 |
| P3 | 5 | 4 |
| P4 | 10 | 10 |
| P5 | 12 | 6 |

■ Consider the following set of process with the length of CPU burst cycle given in milliseconds:


## NEXT



# CPU Scheduling 

## Priority Scheduling



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